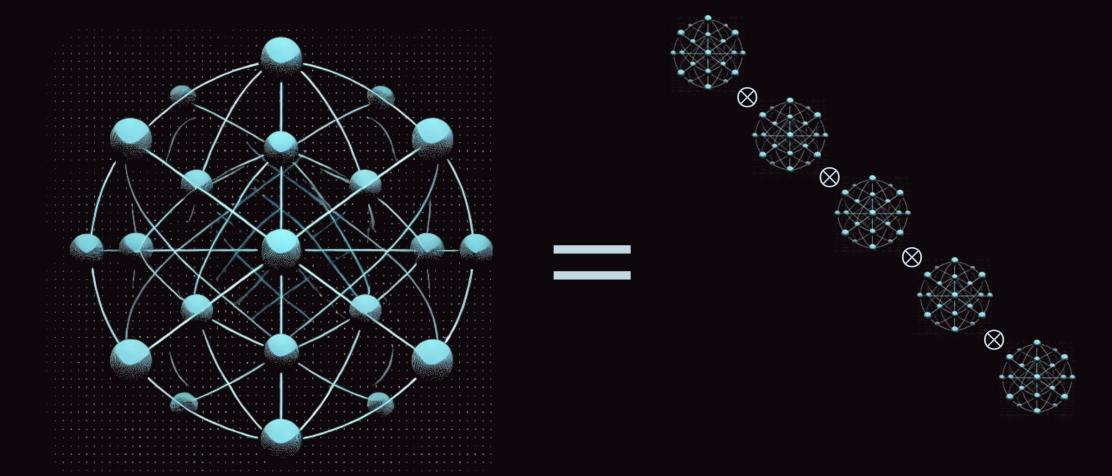
Ainesh Bakshi, John Bostanci, William Kretschmer, Zeph Landau, Jerry Li, Allen Liu, Ryan O'Donnell, and Ewin Tang

Quantum state learning

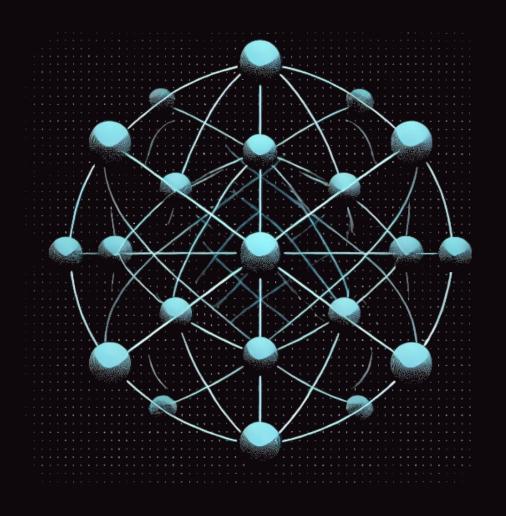
Typical set up:

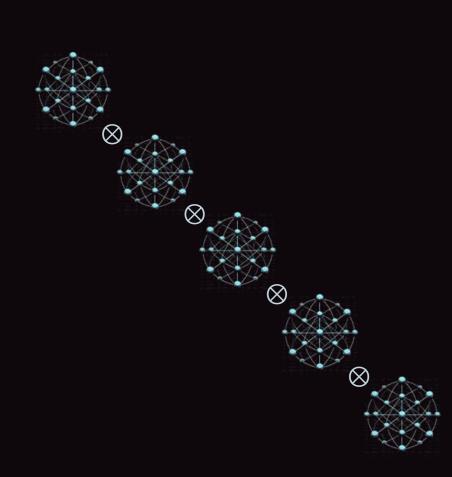


Quantum state learning



Quantum state learning





Task: Learn a classical description of a product state.

Algorithm: Do state tomography on every register and output the tensor product the reduced states.



Consider the following input:

$$\sqrt{1-\epsilon}|0^n\rangle+\sqrt{\epsilon}|+^n\rangle$$

Consider the following input:

$$\sqrt{1-\epsilon}|0^n\rangle+\sqrt{\epsilon}|+^n\rangle$$

Close to all 0's (product state), but every marginal isn't quite $|0\rangle$.

Moral: Even a little bit of misclassification error can change the nature of quantum state learning problems.



Given a model class C and copies of an arbitrary quantum state, output the description of the closest state in C to the state.

$$(C, \rho^{\bigotimes n}) \to \underset{|\psi\rangle \in C}{\operatorname{argmax}} \langle \psi | \rho | \psi \rangle$$

Given a model class C and copies of an arbitrary quantum state, output the description of the closest state in C to the state.

$$\left(C, \rho^{\otimes n}\right) \to |\phi\rangle \in C: \max_{|\psi\rangle \in C} \langle \psi | \rho | \psi\rangle - \langle \phi | \rho | \phi\rangle \leq \epsilon$$

Given a model class C and copies of an arbitrary quantum state, output the description of the closest state in C to the state.

$$\left(C, \rho^{\otimes n}\right) \to |\phi\rangle \in C: \max_{|\psi\rangle \in C} \langle \psi | \rho | \psi\rangle - \langle \phi | \rho | \phi\rangle \le \epsilon$$

If you don't care about runtime, shadow tomography solves this in $O(n \cdot \log^2(|C|) \cdot \epsilon^{-4})$ samples.

Given a model class C and copies of an arbitrary quantum state, output the description of the closest state in C to the state.

$$\left(C, \rho^{\otimes n}\right) \to |\phi\rangle \in C: \max_{|\psi\rangle \in C} \langle \psi | \rho | \psi\rangle - \langle \phi | \rho | \phi\rangle \le \epsilon$$

Surprisingly, few computationally efficient algorithms exist, even for simple families like product states!

Main result: We provide an algorithm for agnostic learning of product states that has sample complexity and time complexity that is

$$\operatorname{poly}(n^{\operatorname{poly}\left(\frac{1}{\epsilon}\right)})$$

We can't use the fact that the input is separable.

We can't use the fact that the input is separable, but Given a state ρ has fidelity at least ½ with some product state.

We can't use the fact that the input is separable, but Given a state ρ has fidelity at least ½ with some product state.

1. The reduced states of ρ have fidelity at least ½ with some product state too.

We can't use the fact that the input is separable, but Given a state ρ has fidelity at least ½ with some product state.

1. The reduced states of ρ have fidelity at least ½ with some product state too.

2. There are at most 2 orthogonal product states that have fidelity larger than $\frac{1}{2}$ with ρ , and all of its reduced states.

These observations motivate the following main loop:

For k from 1 through n:

Given a net $\{\pi_i\}$ for first k registers $(\langle \pi_i | \rho_{[k]} | \pi_i \rangle \ge \frac{1}{2} \text{ and } \langle \pi_i | \pi_j \rangle \approx 0)$, Find a net for k+1 registers.

High level algorithm:

1. Search in a small ball around $|\pi_i\rangle \otimes |\phi_{k+1}\rangle$ (the root candidate).

High level algorithm:

- 1. Search in a small ball around $|\pi_i\rangle \otimes |\phi_{k+1}\rangle$ (the root candidate).
- 2. For a state with high fidelity with ρ .

High level algorithm:

- 1. Search in a small ball around $|\pi_i\rangle \otimes |\phi_{k+1}\rangle$ (the root candidate).
- 2. For a state with high fidelity with ρ .
- 3. That is far from other good candidates we found.

High level algorithm:

- 1. Search in a small ball around $|\pi_i\rangle \otimes |\phi_{k+1}\rangle$ (the root candidate).
- 2. For a state with high fidelity with ρ .
- 3. That is far from other good candidates we found.

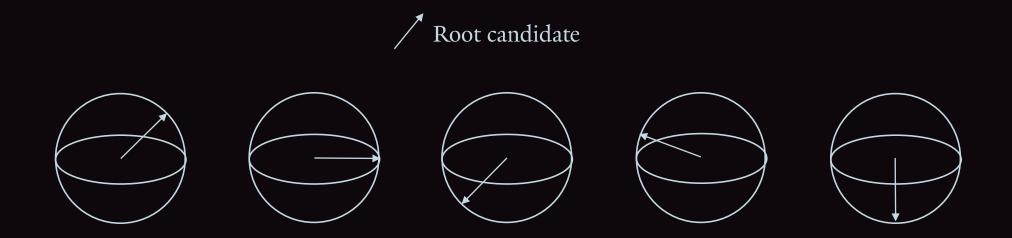
High level algorithm:

- 1. Search in a small ball around $|\pi_i\rangle \otimes |\phi_{k+1}\rangle$ (the root candidate).
- 2. For a state with high fidelity with ρ .
- 3. That is far from other good candidates we found.

The remaining technical challenge will be turning our objective into a low-degree polynomial and then optimizing that polynomial.

Our algorithm:

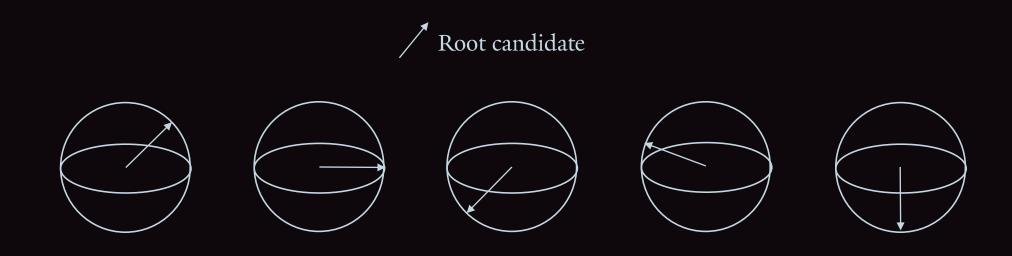
- . Search in a small ball around the root candidate.
- For a state with high fidelity with ρ .
- That is far from other good candidates we found.



- . Search in a small ball around the root candidate.
- 2. For a state with high fidelity with ρ .
- That is far from other good candidates we found.

Learning the closest product state

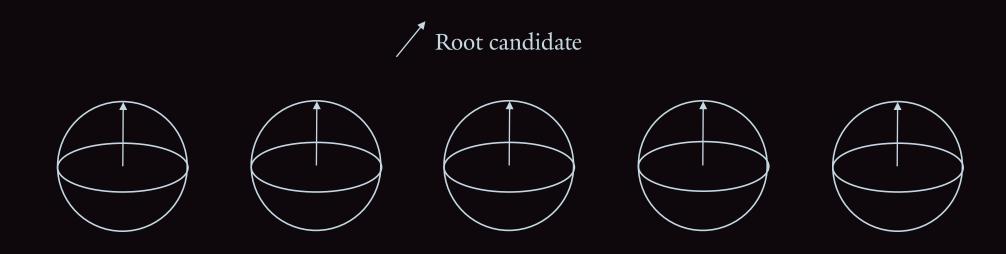
First rotate so that our root candidate is all 0's.



- 1. Search in a small ball around the root candidate.
- 2. For a state with high fidelity with ρ .
- That is far from other good candidates we found.

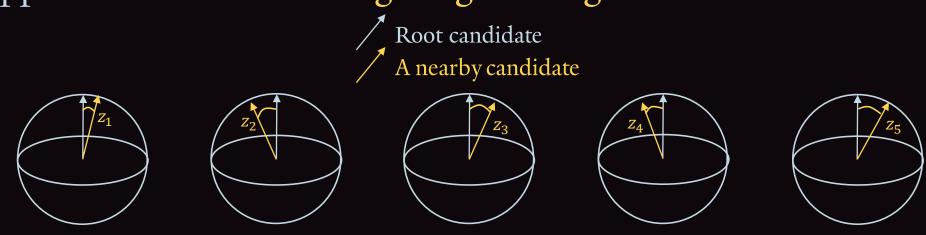
Learning the closest product state

First rotate so that our root candidate is all 0's.



- 1. Search in a small ball around the root candidate.
- 2. For a state with high fidelity with ρ .
 - That is far from other good candidates we found.

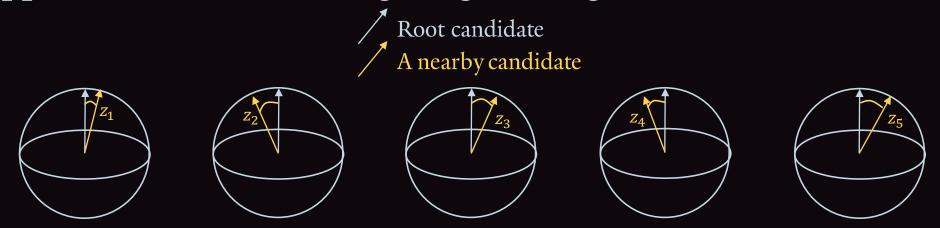
Then the product state ball around $|0^{k+1}\rangle$ will be almost entirely supported on low Hamming weight strings.



- Search in a small ball around the root candidate.
- 2. For a state with high fidelity with ρ .
 - That is far from other good candidates we found.

Learning the closest product state

Then the product state ball around $|0^{k+1}\rangle$ will be almost entirely supported on low Hamming weight strings.



The "quantum part" of the algorithm will be to do tomography on the low-weight restriction of the input ρ .

- 1. Search in a small ball around the root candidate.
- 2. For a state with high fidelity with ρ .
 - That is far from other good candidates we found.

Learning the closest product state

Maximizing fidelity ≈ maximizing the following polynomial

$$\max_{\vec{z} \in \mathbb{C}^{k+1}} \sum_{\substack{x,x' \in \{0,1\}^{k+1} \\ |x|,|x| \leq d}} \langle x | \rho_d | x \rangle (\vec{z}^{*,x}) (\vec{z}^{x'})$$

- Search in a small ball around the root candidate.
- 2. For a state with high fidelity with ρ .
 - That is far from other good candidates we found.

Learning the closest product state

Maximizing fidelity ≈ maximizing the following polynomial

$$\max_{\vec{z} \in \mathbb{C}^{k+1}} \sum_{\substack{x,x' \in \{0,1\}^{k+1} \\ |x|,|x| \le d}} \langle x | \rho_d | x \rangle (\vec{z}^{*,x}) (\vec{z}^{x'})$$
Low Degree!

Improvements in some settings

1. Very high fidelity $(\geq \frac{5}{6})$

Improvements in some settings

1. Very high fidelity $(\geq \frac{5}{6})$

2. Finitely many, far apart, choices per register

Improvements in some settings

- 1. Very high fidelity $(\geq \frac{5}{6})$
- 2. Finitely many, far apart, choices per register
- 3. Polynomial bond-dimension MPS

Even more efficient algorithms?

Our algorithm is only polynomial time when ϵ is a constant. Is there an algorithm that runs in polynomial time when ϵ is small?

Even more efficient algorithms?

Our algorithm is only polynomial time when ϵ is a constant. Is there an algorithm that runs in polynomial time when ϵ is small?

Not unless NP
$$\subseteq$$
 BQP (if $\epsilon \approx \frac{1}{\text{poly}(n)}$).

Even more efficient algorithms?

Our algorithm is only polynomial time when ϵ is a constant. Is there an algorithm that runs in polynomial time when ϵ is small?

Not unless NP
$$\subseteq$$
 BQP (if $\epsilon \approx \frac{1}{\text{poly}(n)}$).

Why? Turns out the connection to tensor optimization goes both ways.